

# Majorizing Cubics on Intervals

Jan de Leeuw

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## Abstract

We illustrate uniform quadratic majorization, sharp quadratic majorization, and sublevel quadratic majorization using the example of a univariate cubic.

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Note: This is a working paper which will be expanded/updated frequently. All suggestions for improvement are welcome. The directory [deleeuwpx.net/pubfolders/cubic](http://deleeuwpx.net/pubfolders/cubic) has a pdf version, the complete Rmd file with all code chunks, the bib file, and the R source code.

# 1 Introduction

Suppose  $\mathcal{I}$  is the closed interval  $[L, U]$ , and  $f : \mathcal{I} \rightarrow \mathbb{R}$ . A function  $g : \mathcal{I} \otimes \mathcal{I} \rightarrow \mathbb{R}$  is a *majorization scheme* for  $f$  on  $\mathcal{I}$  if

- $g(x, x) = f(x)$  for all  $x \in \mathcal{I}$ ,
- $g(x, y) \geq f(x)$  for all  $x, y \in \mathcal{I}$ .

In other words for each  $y$  the global minimum of  $g(x, y) - f(x)$  over  $x \in \mathcal{I}$  is zero, and it is attained at  $y$ . If the functions  $f$  and  $g$  are differentiable and the minimum is attained at an interior point of the interval, we have  $\mathcal{D}_1 g(x, y) = \mathcal{D}f(x)$ . If the functions are in addition twice differentiable we have  $\mathcal{D}_{11} g(x, y) \geq \mathcal{D}^2 f(x)$ .

The majorization conditions are not symmetric in  $x$  and  $y$ , and consequently it sometimes is more clear to write  $g_y(x)$  for  $g(x, y)$ , so that  $g_y : \mathcal{I} \rightarrow \mathbb{R}$ . We say that  $g_y$  majorizes  $f$  on  $\mathcal{I}$  at  $y$ , or with *support point*  $y$ .

A *majorization algorithm* is of the form

$$x^{(k+1)} \in \underset{x \in \mathcal{I}}{\mathbf{argmin}} g(x, x^{(k)})$$

It then follows that

$$f(x^{(k+1)}) \leq g(x^{(k+1)}, x^{(k)}) \leq g(x^{(k)}, x^{(k)}) = f(x^{(k)}), \quad (1)$$

Thus a majorization step decreases the value of the objective function. The chain (1) is called the *sandwich inequality*. In (1) the inequality  $f(x^{(k+1)}) \leq g(x^{(k+1)}, x^{(k)})$  follows from majorization, the inequality  $g(x^{(k+1)}, x^{(k)}) \leq g(x^{(k)}, x^{(k)})$  follows from minimization. This explains why majorization algorithms are also called *MM algorithms* (Lange (2016 (in press))). Using the *MM* label has the advantage that it can also be used for the dual family of minorization-maximization algorithms.

In this note we are interested in *quadratic majorization*, i.e. in majorization functions of the form

$$g(x, y) = f(y) + f'(y)(x - y) + \frac{1}{2}K(x - y)^2, \quad (2)$$

and specifically on quadratic majorizers of a cubic on an closed interval of the real line. If  $K \leq 0$  in (2) the majorization function is concave and attains its minimum at one of endpoints of  $\mathcal{I}$ . If  $K > 0$  then define the *algorithmic map*  $\mathcal{A}(x) = x - f'(x)/K$  and

$$x^{(k+1)} = \begin{cases} L & \text{if } \mathcal{A}(x^{(k)}) < L, \\ \mathcal{A}(x^{(k)}) & \text{if } L \leq \mathcal{A}(x^{(k)}) \leq U, \\ U & \text{if } \mathcal{A}(x^{(k)}) > U. \end{cases}$$

Assuming that the sequence  $x^{(k)}$  converges to a fixed point  $x_\infty$  of  $\mathcal{A}$ , i.e a point with  $\mathcal{D}f(x_\infty) = 0$ , the rate of convergence is

$$\rho(x_\infty) = 1 - \frac{\mathcal{D}^2 f(x_\infty)}{K}.$$

## 2 Majorizing a Cubic

Suppose  $f$  is a non-trivial cubic, with non-zero leading coefficient. The function and its derivatives are

$$\begin{aligned}f(x) &= d + cx + bx^2 + ax^3, \\f'(x) &= c + 2bx + 3ax^2, \\f''(x) &= 2b + 6ax, \\f'''(x) &= 6a.\end{aligned}$$

### 2.1 Uniform Quadratic Majorization

Define

$$K_0 = \max_{L \leq x \leq U} f''(x) = \max(f''(A), f''(B)).$$

Thus

$$K_0 = \begin{cases} 6aU + 2b & \text{if } a > 0, \\ 6aL + 2b & \text{if } a < 0. \end{cases}$$

Our majorization function is the quadratic

$$g(x, y) := f(y) + f'(y)(x - y) + \frac{1}{2}K_0(x - y)^2.$$

The corresponding majorization algorithm is

$$x^{(k+1)} = \mathbf{argmin}_{L \leq x \leq U} g(x, x^{(k)})$$

Note the majorizing function can be a concave quadratic, in which case its minimum is always at one of the endpoints of the interval. Assuming we eventually converge to a value  $L < x_\infty < U$  the convergence rate (or asymptotic error constant) is

$$\rho(x_\infty) := 1 - \frac{f''(x_\infty)}{K_0} = \begin{cases} \frac{6a(U-x_\infty)}{6aU+2b} & \text{if } a > 0, \\ \frac{6a(L-x_\infty)}{6aL+2b} & \text{if } a < 0. \end{cases}$$

Note that this does not depend on the lower limit  $L$  of the interval if  $a > 0$ . Convergence is faster if the upper limit  $U$  happens to be close to  $x_\infty$ , and in fact it can be close to zero. If  $U \rightarrow \infty$  the rate of convergence goes to one.

As a first example, consider  $f(x) = \frac{1}{6}(1 - 2x + x^3)$ . This cubic has roots at -1.6180339887, 0.6180339887, 1. There is a local maximum at  $-\frac{1}{3}\sqrt{6}$  and a local minimum at  $\frac{1}{3}\sqrt{6}$ . The function (red), its first derivative (blue), and its second derivative (green) are in figure 1.

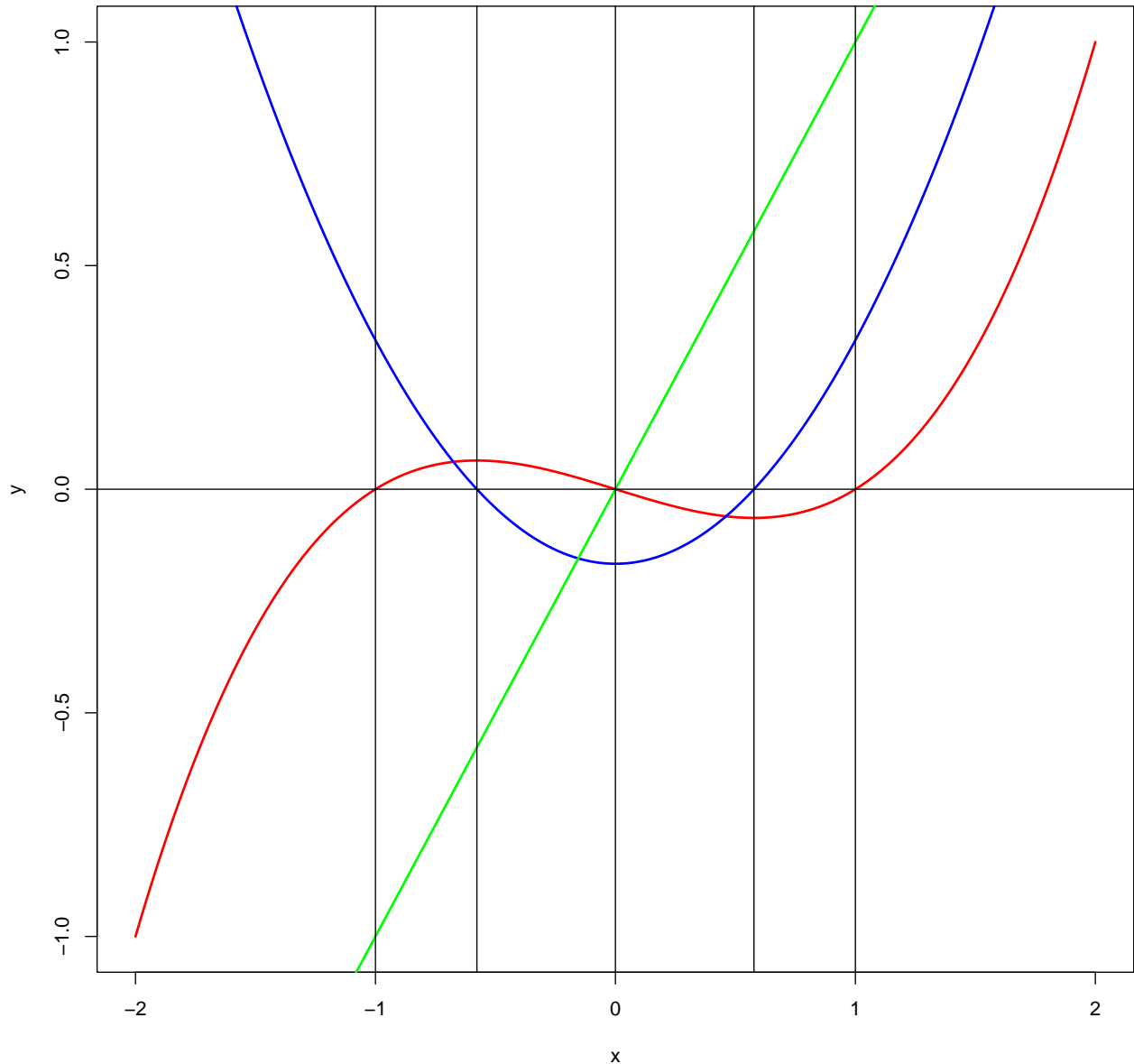


Figure 1: Example Cubic

We will look for a local minimum in the interval  $[-2, 2]$ , stating with initial value 1. Note that  $f''(x) = x$  and thus  $K_0 = B = 2$ . At  $x_\infty = \frac{1}{3}\sqrt{3}$  we have  $\rho(x_\infty) = 1 - \frac{1}{6}\sqrt{3}$ , i.e. approximately 0.7113248654. We report the results of the final iteration.

```
## Iteration: 35 xinit: 1.00000000 xfinal: 0.57735207 rate: 0.71132334
```

If we look for a local minimum in the interval  $[0, 1]$  instead of  $[-2, 2]$ , we get much faster convergence, because the upper bound is now much closer to the solution.

```
## Iteration: 14 xinit: 0.50000000 xfinal: 0.57734974 rate: 0.42265183
```

If we start at a value to the left of  $-\frac{1}{3}\sqrt{3}$  and look for a minimum in  $[-2, 2]$  then the algorithm converges to the boundary at  $-2$ .

```
## Iteration:      3 xinit:  -1.50000000 xfinal:  -2.00000000 rate:  0.00000000
```

Note that alternatively we could have used

$$K_0^+ := \max_{L \leq x \leq U} |f''(x)|$$

in our majorization functions. Since  $f''$  is linear we see that  $|f''|$  is convex, and thus  $K_0^+ = \max(|f''(L)|, |f''(U)|)$ . Using  $K_0^+$  gives a majorization which is generally less precise, but uses majorization functions that are always convex quadratics. Also note that if there is a strict local minimum in  $\mathcal{I}$  then  $K_0 > 0$ , although we can still have  $K_0 < K_0^+$ . Think of  $[-1, .75]$ , for which  $K_0 = .75$  and  $K_0^+ = 1$ .

## 2.2 Sharp Quadratic Majorization

Consider the general representation of a cubic around the point  $y$

$$f(x) = f(y) + f'(y)(x - y) + \frac{1}{2}f''(y)(x - y)^2 + \frac{1}{6}f'''(x - y)^3.$$

For a quadratic function of the form (2) we have

$$g(x, y) - f(x) = \frac{1}{2}(x - y)^2 \left\{ K - f''(y) - \frac{1}{3}f'''(x - y) \right\}. \quad (3)$$

Thus the quadratic function is a majorizer if

$$K \geq \max_{x \in \mathcal{I}} f''(y) + \frac{1}{3}f'''(x - y) \quad (4)$$

which works out to  $K \geq f''(y) + \frac{1}{3} \max(f'''(U - y), f'''(L - y))$ . We get the *sharp quadratic majorization* (De Leeuw and Lange (2009)) by choosing  $K$  equal to its lower bound. The corresponding rate in the cubic case, with positive leading coefficient, is

$$\rho(x_\infty) = \frac{\frac{1}{3}f'''(U - x_\infty)}{f''(x_\infty) + \frac{1}{3}f'''(U - x_\infty)}.$$

If we reanalyze our example with the sharp bound we find faster convergence in the first two computing runs.

```
## Iteration:      17 xinit:  1.00000000 xfinal:  0.57735073 rate:  0.45096147
```

```
## Iteration:       8 xinit:  0.50000000 xfinal:  0.57735010 rate:  0.19615217
```

We also find convergence to the local minimum, and not to the nearby boundary, in the third run.

```
## Iteration:      18 xinit:  -1.50000000 xfinal:  0.57735105 rate:  0.45096119
```

## 2.3 Sublevel Quadratic Majorization

We know that quadratic majorization of a cubic on the whole line is impossible. This is one of the reasons for looking at quadratic majorization on a closed interval, where a continuous second derivative is always bounded. In this section we relax the majorization requirements, using a closed interval that depends on the current solution and becomes smaller if we get closer to a minimum. The resulting majorization method can be thought of as a safeguarded version of Newton's method.

A function  $g : \mathcal{I} \otimes \mathcal{I} \rightarrow \mathbb{R}$  is a *sublevel majorization scheme* for  $f$  on  $\mathcal{I}$  if

- $g(x, x) = f(x)$  for all  $x \in \mathcal{I}$ ,
- $g(x, y) \geq f(x)$  for all  $x, y \in \mathcal{I}$  for which  $g(x, y) \leq g(y, y)$ .

The second part of the definition says that  $g$  majorizes  $f$  on the sublevel set  $\{x \in \mathcal{I} \mid g(x, y) \leq g(y, y)\}$ . If we minimize the sublevel majorization the sandwich inequality (1) is still valid, so we still have monotone convergence of function values.

We quickly specialize this to quadratic majorization functions, suppose  $\mathcal{I}$  is the whole real line, and also require  $K \geq 0$ . If  $g$  is given by (2) then the sublevel set is the interval between  $y$  and  $y - 2f'(y)/K$ . Note that either of the two bounds can be the smaller one. Thus we want inequality

$$K \geq f''(y) + \frac{1}{3}f'''(x - y)$$

on the sublevel set, or equivalently at both endpoints. Thus  $K \geq \max(f''(y), 0)$  and

$$K \geq f''(y) - \frac{2}{3} \frac{f'''f'(y)}{K}.$$

This means we must have

$$K^2 - Kf''(y) + \frac{2}{3}f'''f'(y) \geq 0. \tag{5}$$

Define  $K(y)$  as the smallest  $K \geq \max(f''(y), 0)$  satisfying (5), and we have *sharp sublevel quadratic majorization* (De Leeuw (2006)).

If the quadratic equation corresponding to (5) has no real roots or a single real root, then the inequality (5) is satisfied for all  $K$ , and thus  $K(y) = \max(f''(y), 0)$ . If the equation has two real roots, they are written as  $p(y) \leq q(y)$ . We have  $p(y) + q(y) = f''(y)$ . Thus if  $p(y)$  and  $q(y)$  are non-negative, then  $0 \leq p(y) \leq q(y) \leq f''(y)$ , and consequently  $K(y) = f''(y) \geq 0$ . If  $p(y) \leq 0$  and  $q(y) \geq 0$  then  $q(y) = f''(y) - p(y) \geq f''(y)$  and thus  $K(y) = q(y)$ . If both  $p(y) \leq 0$  and  $q \leq 0$  then  $f''(y) \leq p(y) \leq q(y) \leq 0$ , and thus  $K(y) = 0$  and the sharp sublevel quadratic is linear.

Figure 2 shows for  $y = \frac{1}{2}$  and various values of  $K$  what sublevel majorization looks like. The function is in red, the quadratic sublevel majorization in blue. Note the different lengths of the sublevel intervals. We see that  $K = \frac{1}{2}$  is too small. For  $y = \frac{1}{2}$  the quadratic is  $K^2 = K - \frac{1}{2}K - \frac{1}{36}$ , which has roots  $-0.0504626063$ ,  $0.5504626063$  and thus the sharp sublevel quadratic has  $K$  equal to  $0.5504626063$ .

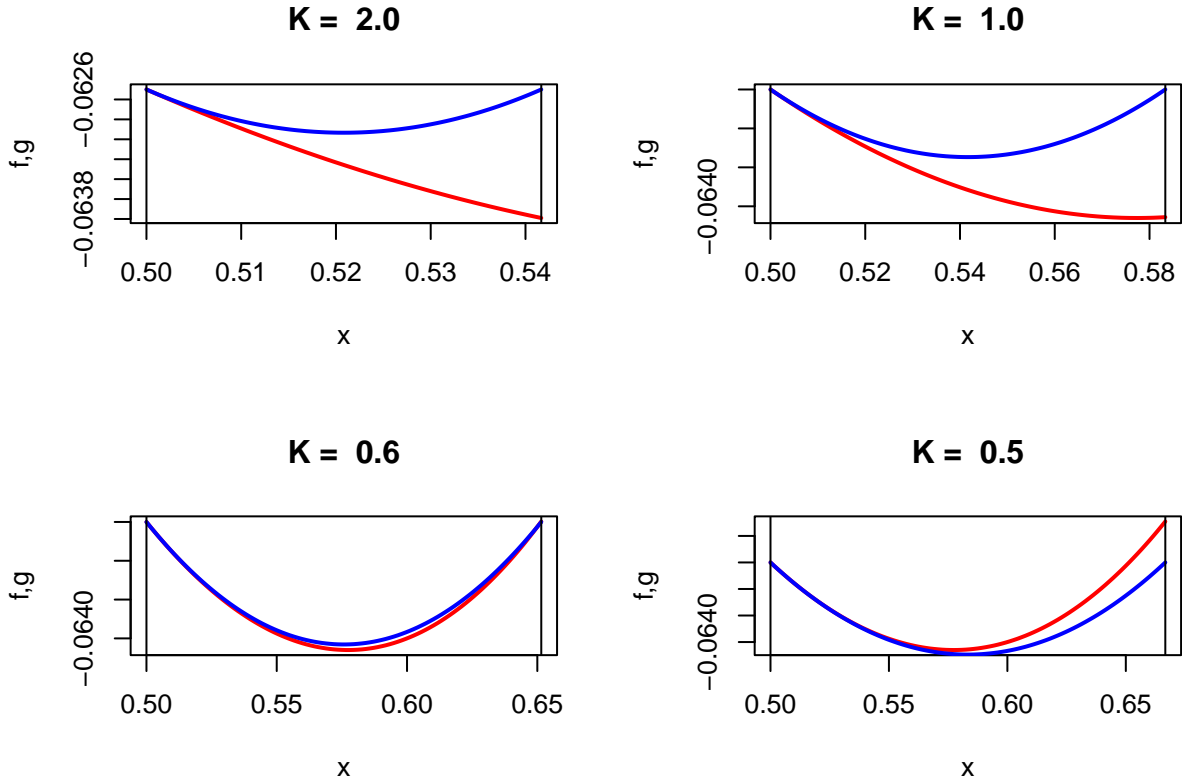


Figure 2: Sublevel Majorization

If we are close to a strict local minimum  $x$  we have  $f'(x) \approx 0$  and  $f''(x) > 0$ . Thus the quadratic will have one root approximately zero and one root approximately equal to  $f''(x)$ , and the iteration is basically a Newton iteration. Thus, at least for cubics, sublevel majorization has quadratic convergence. We illustrate this by analyzing our small example with sublevel majorization, starting from  $x = 1$  and  $x = \frac{1}{2}$ .

```
## Iteration: 1 xold: 1.00000000 xnew: 0.66666667 cnew: 0.33333333 rate:
## Iteration: 2 xold: 0.66666667 xnew: 0.58333333 cnew: 0.08333333 rate:
## Iteration: 3 xold: 0.58333333 xnew: 0.57738095 cnew: 0.00595238 rate:
## Iteration: 4 xold: 0.57738095 xnew: 0.57735027 cnew: 0.00003068 rate:
## Iteration: 5 xold: 0.57735027 xnew: 0.57735027 cnew: 0.00000000 rate:

## Iteration: 1 xold: 0.50000000 xnew: 0.57569391 cnew: 0.07569391 rate:
## Iteration: 2 xold: 0.57569391 xnew: 0.57734948 cnew: 0.00165557 rate:
## Iteration: 3 xold: 0.57734948 xnew: 0.57735027 cnew: 0.00000079 rate:
```

If  $\mathcal{I} = [L, U]$  then our analysis must be modified slightly. We want inequality (4) on the intersection of the sublevel interval and  $[L, U]$ . If the sublevel interval is in  $[L, U]$  our previous analysis applies. Because we always have  $y \in [L, U]$  the other possible intervals are  $[y, U]$  if  $y \leq U \leq y - 2f'(y)/K$  and  $[L, y]$  if  $y - 2f'(y)/K \leq L \leq y$ .

## 3 Appendix: Code

### 3.1 auxiliary.R

```
mprint <- function (x, d = 2, w = 5) {
  print (noquote (formatC (
    x, di = d, wi = w, fo = "f"
  )))
}

cobwebPlotter <-
function (xold,
         func,
         lowx = 0,
         hghx = 1,
         lowy = lowx,
         hghy = hghx,
         eps = 1e-10,
         itmax = 25,
         ...) {
  x <- seq (lowx, hghx, length = 100)
  y <- sapply (x, function (x)
    func (x, ...))
  plot (
    x,
    y,
    xlim = c(lowx, hghx),
    ylim = c(lowy, hghy),
    type = "l",
    col = "RED",
    lwd = 2
  )
  abline (0, 1, col = "BLUE")
  base <- 0
  itel <- 1
  repeat {
    xnew <- func (xold, ...)
    if (itel > 1) {
      lines (matrix(c(xold, xold, base, xnew), 2, 2))
    }
    lines (matrix(c(xold, xnew, xnew, xnew), 2, 2))
    if ((abs (xnew - xold) < eps) || (itel == itmax)) {
      break ()
    }
  }
}
```



```

    }
    base <- xnew
    xold <- xnew
    itel <- itel + 1
  }
}

minQuadratic <- function (a, lw, up) {
  f <- polynomial (a)
  fup <- predict (f, up)
  flw <- predict (f, lw)
  if (a[3] <= 0) {
    if (fup <= flw) return (list (x = up, f = fup))
    if (fup >= flw) return (list (x = lw, f = flw))
  }
  xmn <- - a[2] / (2 * a[3])
  fmn <- predict (f, xmn)
  if (xmn >= up) return (list (x = up, f = fup))
  if (xmn <= lw) return (list (x = lw, f = flw))
  return (list (x = xmn, f = fmn))
}

```

## 3.2 iterate.R

```

myIterator <-
  function (xinit,
            f,
            eps = 1e-6,
            itmax = 100,
            verbose = FALSE,
            final = TRUE,
            ...) {
    xold <- xinit
    cold <- Inf
    itel <- 1
    repeat {
      xnew <- f (xold, ...)
      cnew <- abs (xnew - xold)
      rate <- cnew / cold
      if (verbose)
        cat(
          "Iteration: ",

```

```

formatC (itel, width = 3, format = "d"),
"xold: ",
formatC (
  xold,
  digits = 8,
  width = 12,
  format = "f"
),
"xnew: ",
formatC (
  xnew,
  digits = 8,
  width = 12,
  format = "f"
),
"cnew: ",
formatC (
  cnew,
  digits = 8,
  width = 12,
  format = "f"
),
"rate: ",
formatC (
  rate,
  digits = 8,
  width = 12,
  format = "f"
),
"\n"
)
if ((cnew < eps) || (itel == itmax))
  break
xold <- xnew
cold <- cnew
itel <- itel + 1
}
if (final)
  cat(
    "Iteration: ",
    formatC (itel, width = 3, format = "d"),
    "xinit: ",
    formatC (
      xinit,

```

```

        digits = 8,
        width = 6,
        format = "f"
    ),
    "xfinal: ",
    formatC (
        xnew,
        digits = 8,
        width = 6,
        format = "f"
    ),
    "rate: ",
    formatC (
        rate,
        digits = 8,
        width = 6,
        format = "f"
    ),
    "\n"
)
return (list (
    itel = itel,
    xinit = xinit,
    xfinal = xnew,
    change = cnew,
    rate = rate
))
}

```

```

cubicUQ <- function (x, a, up, lw, sharp = FALSE) {
  f <- polynomial (a)
  g <- deriv (f)
  h <- deriv (g)
  i <- deriv (h)
  if (!sharp) {
    if (a[4] > 0)
      k <- predict (h, up)
    if (a[4] < 0)
      k <- predict (h, lw)
  }
  if (sharp) {
    if (a[4] > 0)
      k <- predict (h, x) + predict (i, x) * (up - x) / 3
    if (a[4] < 0)

```

```

    k <- predict (h, x) + predict (i, x) * (lw - x) / 3
  }
  xmin <- x - predict (g, x) / k
  if ((xmin <= up) && (xmin >= lw))
    return (xmin)
  fup <- predict (f, up)
  flw <- predict (f, lw)
  return (ifelse (fup < flw, up, lw))
}

```

### 3.3 sublevel.R

```

tester <- function (y, k, func, grad) {
  qmaj <- function (x) func (y) + grad (y) * (x - y) + .5 * k * (x - y) ^ 2
  ybnd <- y - 2 * grad (y) / k
  up <- max (y, ybnd)
  lw <- min (y, ybnd)
  x <- seq (lw, up, length = 100)
  s <- paste ("K = ", formatC(k, digits = 1, format= "f"))
  plot (x, func (x), col = "RED", lwd = 2, type = "l", ylab = "f,g", main = s)
  lines (x, qmaj (x), col = "BLUE", lwd = 2)
  abline (v = up)
  abline (v = lw)
}

f <- function (x) log (1 + exp (x))
g <- function (x) exp (x) / (1 + exp (x))

a <- function (x) (x ^ 3 - x) / 6
b <- function (x) (3 * x ^ 2 - 1) / 6

cubicSublevel <- function (y, a, sharp = FALSE) {
  f <- polynomial (a)
  g <- deriv (f)
  h <- deriv (g)
  i <- deriv (h)
  dfy <- predict (g, y)
  dgy <- predict (h, y)
  dhy <- predict (i, y)
  disk <- dgy ^ 2 - 4 * (2 / 3) * dhy * dfy
  if (disk <= 0) k <- max (0, dgy)
  else {

```

```
r <- sort (solve (polynomial (c((2 / 3) * dhy * dfy, -dgy, 1))))
if ((r[1] >= 0) && (r[2] >= 0)) k <- dgy
if ((r[1] <= 0) && (r[2] >= 0)) k <- r[2]
if ((r[1] <= 0) && (r[2] <= 0)) stop ("unbounded")
}
return (y - dfy / k)
}
```

## References

- De Leeuw, J. 2006. “Sharp Local Quadratic Majorization.”
- De Leeuw, J., and K. Lange. 2009. “Sharp Quadratic Majorization in One Dimension.” *Computational Statistics and Data Analysis* 53: 2471–84.
- Lange, K. 2016 (in press). *MM Optimization Algorithms*.