

Constrained Principal Components

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Abstract

In the usual forms of least squares nonlinear principal component analysis observed variables are quantified or transformed to optimize low-rank approximations. Thus NLPCA is linear PCA on optimally scaled variables. In this note we extend the approach by allowing for optimally scaled components.

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Note: This is a working paper which will be expanded/updated frequently. All suggestions for improvement are welcome. The directory deleeuwpx.net/pubfolders/ordered has a pdf version, the bib file, the complete Rmd file with the code chunks, and the R and C source code. Thanks to Yoshio for some useful comments.

1 Loss Function

The problem studied in this paper is minimization of the least squares loss function

$$\sigma(X, B) = \mathbf{SSQ}(Y - X\underline{B}) \tag{1}$$

over X and B . Here $\mathbf{SSQ}()$ stands for the sum of squares, i.e. the square of the Frobenius norm. We use \underline{B} for the matrix transpose of B . The (partial) unknowns are X , an $n \times p$ matrix of *components*, and B , an $m \times p$ matrix of *loadings*. This is just least squares low-rank approximation. It is typically solved by directly using the singular value decomposition of Y , or by using some version of alternating least squares (iterate finding the optimal X for given B and the optimal B for given X).

In order to at least partially identify X and B we require, without loss of generality, that $\underline{X}X = I$, i.e. the components are orthonormal. This identifies X and B up to a rotation. All of this is completely standard. In this paper we introduce additional non-trivial constraints on the components. As an example, we could require for the first component

$$x_{11} \leq x_{21} \leq \dots \leq x_{n1}, \tag{2}$$

but general partial orders and other cone or subspace constraints may also be useful. The important thing is that the constraints are imposed on each component separately.

We write Ω_s for the set of all matrices satisfying the constraint for component s , which means we require $X \in \Omega$, with

$$\Omega = \Omega_1 \otimes \dots \otimes \Omega_p.$$

This form of nonlinear principal component analysis is different from the more familiar form in which the columns of Y are transformed nonlinearly to optimize (1). See De Leeuw (2006) or De Leeuw (2014) for discussion and references. It is more closely related to the forms of constrained principal component analysis discussed in great detail by Takane (2014), and specifically to the form that allows separate constraints for separate dimensions (Takane, Kiers, and De Leeuw (1995)). In constrained principal component analysis, however, the emphasis is primarily on linear or subspace constraints on the components.

2 Alternating Least Squares

If we apply block relaxation to a least squares loss function we obtain an alternating least squares algorithm (De Leeuw (1994), De Leeuw (2020)). Alternating least squares algorithms as a general class of algorithms useful in data analysis were introduced by De Leeuw (1968), and then applied systematically in the ALSOS project, starting with De Leeuw, Young, and Takane (1976).

0: Start with $k = 0$ and $X^{(0)} \in \Omega$.

1: $B^{(k)} \in \mathbf{argmin}_{B \in \mathbb{R}^{m \times p}} \mathbf{SSQ}(Y - X^{(k)}\underline{B})$

2: $X^{(k+1)} \in \mathbf{argmin}_{X \in \Omega} \mathbf{SSQ}(Y - X\underline{B}^k)$

3: If there is convergence, stop, otherwise $k \leftarrow k + 1$ and go back to step k1.

Step 1 is a straightforward linear least squares problem, since we do not impose any constraints on B . Step 2 is more complicated, and we'll discuss it in detail in the next section. Convergence can be defined in terms of the decreasing sequence of loss function values, or in terms of changes in X and B from one iteration to the next.

3 Majorization

Consider the problem of minimizing $\mathbf{SSQ}(Y - X\underline{B})$ over $X \in \mathcal{K}$. Although the constraints are defined for each column of X separately, the matrix B combines columns and thus complicates the overall minimization problem. In order to solve the problem we use the majorization method, introduced by De Leeuw (1994), Heiser (1995), and Lange, Hunter, and Yang (2000), to separate the columns again. Majorization methods are discussed systematically in Lange (2016) and De Leeuw (2021). In our overall algorithm this means we need an iterative majorization algorithm in step 2 of each cycle of the alternating least squares algorithm.

For our majorization we choose a positive semi-definite diagonal matrix W such that $W \succeq \underline{B}B$, i.e. such that $W - \underline{B}B$ is positive semi-definite. We can use, for example, $W = \lambda I$, with λ the largest eigenvalue of $\underline{B}B$, or we can choose $W = p \mathbf{diag}(\underline{B}B)$, with p the number of columns of X and B . Of course if $\underline{B}B$ already is diagonal we can choose $W = \underline{B}B$.

Suppose $Z \in \mathcal{K}$, and in the majorization iteration we want improve Z . By defining

$$U = Z + (Y - Z\underline{B})BW^{-1}$$

and completing the square, we find the majorization inequality

$$\mathbf{SSQ}(Y - X\underline{B}) \leq \mathbf{SSQ}(Y - Z\underline{B}) + \mathbf{tr}(X - U)W(X - U) - \mathbf{tr}UWU,$$

with equality if $Z = X$.

In the majorization algorithm we iteratively minimize

$$\mathbf{tr}(X - U)W(X - U) = \sum_{s=1}^p w_s \mathbf{SSQ}(x_{*s} - u_{*s})$$

over $x_{*s} \in \mathcal{K}_s$, which can obviously be done for each s separately. The update of column x_{*s} of X is the metric projection of column u_{*s} of U on Ω_s . For constraint (2), for example, we apply isotone regression to u_{*s} .

4 Example: Linear Constraints

Our first example is in the spirit of the DCDD method in Takane (2014). Each column of X is constrained linearly by $x_{*s} = G_s \alpha_s$. Thus x_{*s} must be in the subspace spanned by the

columns of G_s . The G_s can be design matrices, indicator matrices, bases of polynomials or splines, and whatever else.

Here is a simple example.

```
g1 <- matrix (0, 16, 4)
g1[1:4, 1] <- 1
g1[5:8, 2] <- 1
g1[9:12, 3] <- 1
g1[13:16, 4] <- 1
g1 <- standardize (center (g1))
g2 <- rbind (diag (4), diag (4), diag (4), diag (4))
g2 <- standardize (center (g2))
g <- list (g1, g2)
set.seed (12345)
y <- standardize (center (mnorm(16, 5)))
```

The `linRes` function approximates Y using the two constrained dimensions code in the list with matrices G_1 and G_2 . Note that in the current implementation we only do a single majorization iteration to update the columns of X before we update B . A more flexible strategy would be to allow for more than one inner majorization iterations.

```
h <- linRes (y, g)
```

```
## itel      1 fold      4.6627879883 fnew      4.6085187514
## itel      2 fold      4.6085187514 fnew      4.5490565074
## itel      3 fold      4.5490565074 fnew      4.4865061504
## itel      4 fold      4.4865061504 fnew      4.4189832984
## itel      5 fold      4.4189832984 fnew      4.3676298613
## itel      6 fold      4.3676298613 fnew      4.3396567038
## itel      7 fold      4.3396567038 fnew      4.3279401130
## itel      8 fold      4.3279401130 fnew      4.3238626172
## itel      9 fold      4.3238626172 fnew      4.3226023894
## itel     10 fold      4.3226023894 fnew      4.3222239136
## itel     11 fold      4.3222239136 fnew      4.3220998929
## itel     12 fold      4.3220998929 fnew      4.3220505942
## itel     13 fold      4.3220505942 fnew      4.3220267104
## itel     14 fold      4.3220267104 fnew      4.3220135396
## itel     15 fold      4.3220135396 fnew      4.3220058435
## itel     16 fold      4.3220058435 fnew      4.3220012301
## itel     17 fold      4.3220012301 fnew      4.3219984440
## itel     18 fold      4.3219984440 fnew      4.3219967506
## itel     19 fold      4.3219967506 fnew      4.3219957194
## itel     20 fold      4.3219957194 fnew      4.3219950873
```

## itel	21 fold	4.3219950873	fnew	4.3219946980
## itel	22 fold	4.3219946980	fnew	4.3219944556
## itel	23 fold	4.3219944556	fnew	4.3219943028
## itel	24 fold	4.3219943028	fnew	4.3219942048
## itel	25 fold	4.3219942048	fnew	4.3219941406
## itel	26 fold	4.3219941406	fnew	4.3219940973
## itel	27 fold	4.3219940973	fnew	4.3219940671
## itel	28 fold	4.3219940671	fnew	4.3219940454
## itel	29 fold	4.3219940454	fnew	4.3219940292
## itel	30 fold	4.3219940292	fnew	4.3219940166
## itel	31 fold	4.3219940166	fnew	4.3219940066
## itel	32 fold	4.3219940066	fnew	4.3219939985
## itel	33 fold	4.3219939985	fnew	4.3219939917
## itel	34 fold	4.3219939917	fnew	4.3219939860
## itel	35 fold	4.3219939860	fnew	4.3219939810
## itel	36 fold	4.3219939810	fnew	4.3219939768
## itel	37 fold	4.3219939768	fnew	4.3219939731
## itel	38 fold	4.3219939731	fnew	4.3219939699
## itel	39 fold	4.3219939699	fnew	4.3219939671
## itel	40 fold	4.3219939671	fnew	4.3219939646
## itel	41 fold	4.3219939646	fnew	4.3219939625
## itel	42 fold	4.3219939625	fnew	4.3219939606
## itel	43 fold	4.3219939606	fnew	4.3219939589
## itel	44 fold	4.3219939589	fnew	4.3219939575
## itel	45 fold	4.3219939575	fnew	4.3219939562
## itel	46 fold	4.3219939562	fnew	4.3219939550
## itel	47 fold	4.3219939550	fnew	4.3219939540
## itel	48 fold	4.3219939540	fnew	4.3219939532
## itel	49 fold	4.3219939532	fnew	4.3219939524
## itel	50 fold	4.3219939524	fnew	4.3219939517
## itel	51 fold	4.3219939517	fnew	4.3219939511
## itel	52 fold	4.3219939511	fnew	4.3219939506
## itel	53 fold	4.3219939506	fnew	4.3219939501
## itel	54 fold	4.3219939501	fnew	4.3219939497
## itel	55 fold	4.3219939497	fnew	4.3219939494
## itel	56 fold	4.3219939494	fnew	4.3219939490
## itel	57 fold	4.3219939490	fnew	4.3219939488
## itel	58 fold	4.3219939488	fnew	4.3219939485
## itel	59 fold	4.3219939485	fnew	4.3219939483
## itel	60 fold	4.3219939483	fnew	4.3219939481
## itel	61 fold	4.3219939481	fnew	4.3219939479
## itel	62 fold	4.3219939479	fnew	4.3219939478
## itel	63 fold	4.3219939478	fnew	4.3219939477
## itel	64 fold	4.3219939477	fnew	4.3219939475
## itel	65 fold	4.3219939475	fnew	4.3219939474

```
## itel      66 fold      4.3219939474 fnew      4.3219939474
```

5 Example: Multiple Ordinal Variables

In the nonlinear multivariate analysis system of Gifi (1990) (see also Michailidis and De Leeuw (1998)) transformations can be nominal, ordinal, or numerical, and quantifications can be single or multiple. Multiple ordinal transformations have not really been implemented because their definition has never been entirely obvious.

A definition that has been used before is to require that the first column x_{*1} is isotone, while the remaining $p-1$ columns are arbitrary, but orthogonal to the first. Thus we have $\underline{X}\underline{X} = I$ with an isotone first column. It is easy to see that the orthonormality constraint is actually just for identification purposes, it does not enter into the majorization iterations. In fact, the algorithm in the function `multOrd` does not impose orthonormality. It uses the modified Gram-Schmidt method to transform to orthonormality after convergence, using the fact that modified Gram-Schmidt (without pivoting) merely normalizes the first column, which does not disturb the order relations. Clearly the same trick can be used for other cone constraints.

It is interesting to observe that requiring isotonicity of the first column of X is actually equivalent to requiring that an isotone vector exists in the column space of X . Because if such a vector exists we can use the indeterminacy in the product $X\underline{B}$ to move that vector to the first column of X .

For our numerical example we use the same Y as before.

```
set.seed (12345)
y <- standardize (center (mnorm(16, 5)))
```

The program `multOrd` require the first column of X to be increasing. It uses Gram-Schmidt from De Leeuw (2015) and isotone regression from De Leeuw (2016).

```
h <- multOrd (y, 2)
```

```
## itel      1 fold      2.9238552791 fnew      2.3439684622
## itel      2 fold      2.3439684622 fnew      2.0718818606
## itel      3 fold      2.0718818606 fnew      2.0413535991
## itel      4 fold      2.0413535991 fnew      2.0263303595
## itel      5 fold      2.0263303595 fnew      2.0212415190
## itel      6 fold      2.0212415190 fnew      2.0177619980
## itel      7 fold      2.0177619980 fnew      2.0149648432
## itel      8 fold      2.0149648432 fnew      2.0127152559
## itel      9 fold      2.0127152559 fnew      2.0108485613
## itel     10 fold      2.0108485613 fnew      2.0093161023
## itel     11 fold      2.0093161023 fnew      2.0080300445
```

## itel	12 fold	2.0080300445	fnew	2.0069592899
## itel	13 fold	2.0069592899	fnew	2.0060533653
## itel	14 fold	2.0060533653	fnew	2.0052907927
## itel	15 fold	2.0052907927	fnew	2.0046414247
## itel	16 fold	2.0046414247	fnew	2.0040901592
## itel	17 fold	2.0040901592	fnew	2.0036182740
## itel	18 fold	2.0036182740	fnew	2.0032150474
## itel	19 fold	2.0032150474	fnew	2.0028684164
## itel	20 fold	2.0028684164	fnew	2.0025707084
## itel	21 fold	2.0025707084	fnew	2.0023138952
## itel	22 fold	2.0023138952	fnew	2.0020924447
## itel	23 fold	2.0020924447	fnew	2.0019008666
## itel	24 fold	2.0019008666	fnew	2.0017351430
## itel	25 fold	2.0017351430	fnew	2.0015914345
## itel	26 fold	2.0015914345	fnew	2.0014668028
## itel	27 fold	2.0014668028	fnew	2.0013585141
## itel	28 fold	2.0013585141	fnew	2.0012644056
## itel	29 fold	2.0012644056	fnew	2.0011825024
## itel	30 fold	2.0011825024	fnew	2.0011112032
## itel	31 fold	2.0011112032	fnew	2.0010490649
## itel	32 fold	2.0010490649	fnew	2.0009948954
## itel	33 fold	2.0009948954	fnew	2.0009476303
## itel	34 fold	2.0009476303	fnew	2.0009063783
## itel	35 fold	2.0009063783	fnew	2.0008703482
## itel	36 fold	2.0008703482	fnew	2.0008388706
## itel	37 fold	2.0008388706	fnew	2.0008113543
## itel	38 fold	2.0008113543	fnew	2.0007872946
## itel	39 fold	2.0007872946	fnew	2.0007662472
## itel	40 fold	2.0007662472	fnew	2.0007478306
## itel	41 fold	2.0007478306	fnew	2.0007317096
## itel	42 fold	2.0007317096	fnew	2.0007175951
## itel	43 fold	2.0007175951	fnew	2.0007052332
## itel	44 fold	2.0007052332	fnew	2.0006944042
## itel	45 fold	2.0006944042	fnew	2.0006849153
## itel	46 fold	2.0006849153	fnew	2.0006765992
## itel	47 fold	2.0006765992	fnew	2.0006693093
## itel	48 fold	2.0006693093	fnew	2.0006629180
## itel	49 fold	2.0006629180	fnew	2.0006573133
## itel	50 fold	2.0006573133	fnew	2.0006523977
## itel	51 fold	2.0006523977	fnew	2.0006480858
## itel	52 fold	2.0006480858	fnew	2.0006443030
## itel	53 fold	2.0006443030	fnew	2.0006409838
## itel	54 fold	2.0006409838	fnew	2.0006380711
## itel	55 fold	2.0006380711	fnew	2.0006355148
## itel	56 fold	2.0006355148	fnew	2.0006332710

## itel	57 fold	2.0006332710	fnew	2.0006313014
## itel	58 fold	2.0006313014	fnew	2.0006295723
## itel	59 fold	2.0006295723	fnew	2.0006280541
## itel	60 fold	2.0006280541	fnew	2.0006267211
## itel	61 fold	2.0006267211	fnew	2.0006255505
## itel	62 fold	2.0006255505	fnew	2.0006245225
## itel	63 fold	2.0006245225	fnew	2.0006236196
## itel	64 fold	2.0006236196	fnew	2.0006228267
## itel	65 fold	2.0006228267	fnew	2.0006221302
## itel	66 fold	2.0006221302	fnew	2.0006215183
## itel	67 fold	2.0006215183	fnew	2.0006209809
## itel	68 fold	2.0006209809	fnew	2.0006205087
## itel	69 fold	2.0006205087	fnew	2.0006200939
## itel	70 fold	2.0006200939	fnew	2.0006197295
## itel	71 fold	2.0006197295	fnew	2.0006194093
## itel	72 fold	2.0006194093	fnew	2.0006191279
## itel	73 fold	2.0006191279	fnew	2.0006188807
## itel	74 fold	2.0006188807	fnew	2.0006186635
## itel	75 fold	2.0006186635	fnew	2.0006184726
## itel	76 fold	2.0006184726	fnew	2.0006183048
## itel	77 fold	2.0006183048	fnew	2.0006181574
## itel	78 fold	2.0006181574	fnew	2.0006180279
## itel	79 fold	2.0006180279	fnew	2.0006179140
## itel	80 fold	2.0006179140	fnew	2.0006178140
## itel	81 fold	2.0006178140	fnew	2.0006177260
## itel	82 fold	2.0006177260	fnew	2.0006176487
## itel	83 fold	2.0006176487	fnew	2.0006175808
## itel	84 fold	2.0006175808	fnew	2.0006175211
## itel	85 fold	2.0006175211	fnew	2.0006174686
## itel	86 fold	2.0006174686	fnew	2.0006174225
## itel	87 fold	2.0006174225	fnew	2.0006173820
## itel	88 fold	2.0006173820	fnew	2.0006173463
## itel	89 fold	2.0006173463	fnew	2.0006173150
## itel	90 fold	2.0006173150	fnew	2.0006172875
## itel	91 fold	2.0006172875	fnew	2.0006172633
## itel	92 fold	2.0006172633	fnew	2.0006172420
## itel	93 fold	2.0006172420	fnew	2.0006172233
## itel	94 fold	2.0006172233	fnew	2.0006172069
## itel	95 fold	2.0006172069	fnew	2.0006171924
## itel	96 fold	2.0006171924	fnew	2.0006171797
## itel	97 fold	2.0006171797	fnew	2.0006171686
## itel	98 fold	2.0006171686	fnew	2.0006171588
## itel	99 fold	2.0006171588	fnew	2.0006171501
## itel	100 fold	2.0006171501	fnew	2.0006171425
## itel	101 fold	2.0006171425	fnew	2.0006171359

## itel	102	fold	2.0006171359	fnew	2.0006171300
## itel	103	fold	2.0006171300	fnew	2.0006171249
## itel	104	fold	2.0006171249	fnew	2.0006171203
## itel	105	fold	2.0006171203	fnew	2.0006171164
## itel	106	fold	2.0006171164	fnew	2.0006171129
## itel	107	fold	2.0006171129	fnew	2.0006171098
## itel	108	fold	2.0006171098	fnew	2.0006171071
## itel	109	fold	2.0006171071	fnew	2.0006171047
## itel	110	fold	2.0006171047	fnew	2.0006171026
## itel	111	fold	2.0006171026	fnew	2.0006171008
## itel	112	fold	2.0006171008	fnew	2.0006170992
## itel	113	fold	2.0006170992	fnew	2.0006170977
## itel	114	fold	2.0006170977	fnew	2.0006170965
## itel	115	fold	2.0006170965	fnew	2.0006170954
## itel	116	fold	2.0006170954	fnew	2.0006170944
## itel	117	fold	2.0006170944	fnew	2.0006170936
## itel	118	fold	2.0006170936	fnew	2.0006170928
## itel	119	fold	2.0006170928	fnew	2.0006170922
## itel	120	fold	2.0006170922	fnew	2.0006170916
## itel	121	fold	2.0006170916	fnew	2.0006170911
## itel	122	fold	2.0006170911	fnew	2.0006170906
## itel	123	fold	2.0006170906	fnew	2.0006170903
## itel	124	fold	2.0006170903	fnew	2.0006170899
## itel	125	fold	2.0006170899	fnew	2.0006170896
## itel	126	fold	2.0006170896	fnew	2.0006170893
## itel	127	fold	2.0006170893	fnew	2.0006170891
## itel	128	fold	2.0006170891	fnew	2.0006170889
## itel	129	fold	2.0006170889	fnew	2.0006170887
## itel	130	fold	2.0006170887	fnew	2.0006170886
## itel	131	fold	2.0006170886	fnew	2.0006170884
## itel	132	fold	2.0006170884	fnew	2.0006170883
## itel	133	fold	2.0006170883	fnew	2.0006170882
## itel	134	fold	2.0006170882	fnew	2.0006170881

6 Appendix: Weights

In least squares majorization with non-diagonal positive definite weight matrix C we often want to find a diagonal matrix such that $D \succeq C$ and D is, in some sense, as small as possible (cf Groenen, Giaquinto, and Kiers (2003)). This problem is similar, but not identical, to minimum trace factor analysis (MTFA). In MTFA we find a diagonal $D \succeq 0$ such that $D \preceq C$ and D is as large as possible. Typically “large” is defined as maximizing the trace or some other linear function of the diagonal elements. The majorization problem is different, not so much because “as large as possible” is replaced by “as small as possible”, but more

so because $D \succeq C$ implies that $D \succeq 0$.

For our majorization we could also use the trace, so we want to minimize $\text{tr } D$ over $D \succeq C$. As in MTFA, this is a convex programming problem. It can be solved in many ways (see, for example, Jamshidian and Bentler (1998)), but each of them involves a non-trivial computational effort. In the body of the paper we have mentioned two more simple, and more easily computable, choices for D . The first is $D = \lambda(C)I$, where $\lambda(C)$ is the largest eigenvalue of C . Of course if C is large computing the largest eigenvalue is not entirely trivial either. The second choice, which is actually used in the R programs in the appendix, is $D = p \text{diag}(C)$, where p is the order of C and D . This is trivial to compute, but the corresponding D may not be very good (too large). It uses the trace as an upper bound for the largest eigenvalue of the correlation matrix corresponding with C , and there are much better bounds.

This suggests using better bounds for the maximum eigenvalue, which are relatively easy to compute. We use the result that $\lambda(C) \leq \|C\|$, where $\|C\|$ is any matrix norm. We could use the Frobenius norm $\sqrt{\text{tr } C^2}$ or the norm $\max_{s=1}^p \sum_{t=1}^p |c_{st}|$.

Let's illustrate this with the 2×2 matrix

$$\begin{bmatrix} 2 & -1 \\ -1 & 3 \end{bmatrix}$$

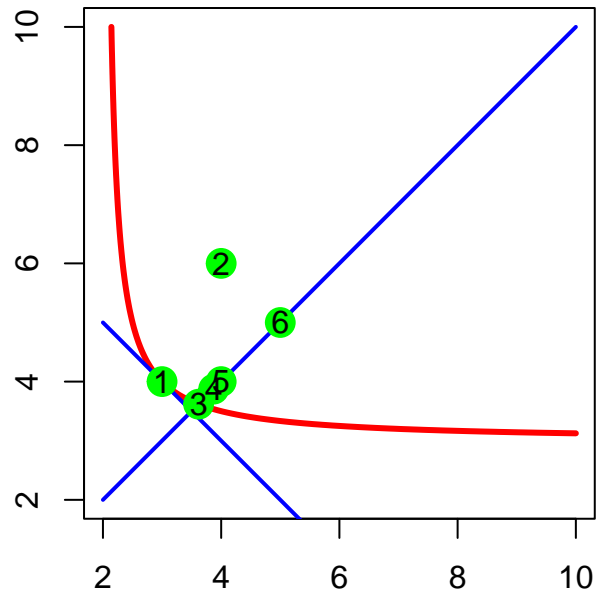
In the figure below all points in the convex area above the red branch from the hyperbola $(x-2)(y-3) = 1$ are D for which $D \succeq C$. The trace of D is minimized at $(3, 4)$, where the blue line $x + y = 7$ is tangent to the hyperbola. This is the point labeled 1. Point labeled 2 is $p \text{diag}(C)$, which is $(4, 6)$. The other four points are on the other blue line $x = y$. Point 3 is at $x = y = (5 + \sqrt{5})/2$, which corresponds with the largest eigenvalue of C . Point 4 has $x = y = \sqrt{15}$, which is the Frobenius bound, and point 5 has $x = y = 4$, which is the maximum row absolute sum norm. Finally point 6 uses the trace as the eigenvalue bound, which gives $x = y = 5$.

```
par(pty="s")
f<-function (x,y) (x-2)*(y-3)-1
y<-seq(2,10,length=100)
x<-seq(2,10,length=100)
z<-outer(x,y,f)
contour(x=x,y=y,z=z,level=0,drawlabels=FALSE,lwd=3,col="RED")
lines(x,7-x,col="BLUE",lwd=2)
lines(x, x, col = "BLUE", lwd = 2)
lbd <- (5+sqrt(5))/2
sbd = sqrt(15)
points (lbd, lbd, col ="GREEN", cex = 2, pch = 19)
points (4, 6, col ="GREEN", cex = 2, pch = 19)
points (3, 4, col ="GREEN", cex = 2, pch = 19)
points (4, 4, col = "GREEN", cex = 2, pch = 19)
points (sbd, sbd, col = "GREEN", cex = 2, pch = 19)
```

```

points (5, 5, col = "GREEN", cex = 2, pch = 19)
text(3,4, labels = "1")
text(4,6, labels = "2")
text(lbd,lbd, labels = "3")
text(sbd,sbd, labels = "4")
text(4,4, labels = "5")
text(5,5, labels = "6")

```



7 Appendix: Code

7.1 linRes.R

```

linRes <-
  function (y,
            g,
            bound = "M",
            itmax = 1000,
            eps = 1e-10,
            verbose = TRUE,
            center = FALSE,
            standardize = FALSE) {
    if (center) {
      y <- center (y)
    }
    if (standardize) {

```

```

y <- standardize (y)
}
p <- length (g)
s <- 1
xold <- NULL
while (s <= p) {
  xold <- cbind (xold, g[[s]] %*% 1:ncol (g[[s]]))
  s <- s + 1
}
bold <- t (lm.fit (xold, y)$coefficients)
roid <- y - tcrossprod (xold, bold)
fold <- ssq (roid)
itel <- 1
repeat {
  bnew <- t (lm.fit (xold, y)$coefficients)
  cnew <- crossprod (bnew)
  if (bound == "M") {
    e <- max (rowSums (abs (cnew)))
    w <- diag (p) / e
  }
  if (bound == "E") {
    e <- max (eigen (cnew)$values)
    w <- diag (p) / e
  }
  if (bound == "F") {
    e <- sqrt (sum (cnew ^ 2))
    w <- diag (p) / e
  }
  if (bound == "D") {
    w <- diag (1 / (p * diag (cnew)))
  }
  u <- xold + roid %*% bnew %*% w
  xnew <- NULL
  s <- 1
  while (s <= p) {
    xnew <- cbind (xnew, lm.fit (g[[s]], u[, s])$fitted.values)
    s <- s + 1
  }
  rnew <- y - tcrossprod (xnew, bnew)
  fnew <- ssq (rnew)
  if (verbose) {
    cat (
      "itel ",
      formatC (itel, width = 4, format = "d"),

```

```

    "fold ",
    formatC (
      fold,
      width = 15,
      digits = 10,
      format = "f"
    ),
    "fnew ",
    formatC (
      fnew,
      width = 15,
      digits = 10,
      format = "f"
    ),
    "\n"
  )
}
if ((itel == itmax) || ((fold - fnew) < eps))
  break
itel <- itel + 1
fold <- fnew
xold <- xnew
bold <- bnew
rold <- rnew
}
return (list (
  x = xnew,
  b = bnew,
  r = rnew,
  f = fnew
))
})
}

```

7.2 multOrd.R

```

source ("jbkPava.R")
source ("gs.R")

multOrd <-
  function (y,
            p,
            bound = "M",

```

```

        itmax = 1000,
        eps = 1e-10,
        verbose = TRUE,
        center = FALSE,
        standardize = FALSE) {
set.seed (12345)
if (center) {
  y <- center (y)
}
if (standardize) {
  y <- standardize (y)
}
n <- nrow (y)
xold <- standardize (center (cbind (1:n, mnorm (n, p - 1))))
bold <- t (lm.fit (xold, y)$coefficients)
rold <- y - tcrossprod (xold, bold)
fold <- ssq (rold)
itel <- 1
repeat {
  bnew <- t (lm.fit (xold, y)$coefficients)
  cnew <- crossprod (bnew)
  if (bound == "M") {
    e <- max (rowSums (abs (cnew)))
    w <- diag (p) / e
  }
  if (bound == "E") {
    e <- max (eigen (cnew)$values)
    w <- diag (p) / e
  }
  if (bound == "F") {
    e <- sqrt (sum (cnew ^ 2))
    w <- diag (p) / e
  }
  if (bound == "D") {
    w <- diag (1 / (p * diag (cnew)))
  }
  xnew <- xold + rold %*% bnew %*% w
  xnew[, 1] <- jbkPava (xnew[, 1])
  rnew <- y - tcrossprod (xnew, bnew)
  fnew <- ssq (rnew)
  if (verbose) {
    cat (
      "itel ",
      formatC (itel, width = 4, format = "d"),

```

```

    "fold ",
    formatC (
      fold,
      width = 15,
      digits = 10,
      format = "f"
    ),
    "fnew ",
    formatC (
      fnew,
      width = 15,
      digits = 10,
      format = "f"
    ),
    "\n"
  )
}
if ((itel == itmax) || ((fold - fnew) < eps))
  break
itel <- itel + 1
fold <- fnew
xold <- xnew
bold <- bnew
rold <- rnew
}
z <- gsrc(xnew)
return (list (x = z$q, b = tcrossprod (bnew, z$r), r = rnew, f = fnew))
}

```

7.3 auxiliary.R

```

mnorm <- function (n, p) {
  return (matrix (rnorm (n * p), n, p))
}

center <- function (y) {
  return (apply (y, 2, function (x)
    x - mean (x)))
}

standardize <- function (y) {
  return (apply (y, 2, function (x)

```

```

    x / norm (x))
}

ssq <- function (x) {
  return (sum (x ^ 2))
}

norm <- function (x) {
  return (sqrt (ssq (x)))
}

```

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